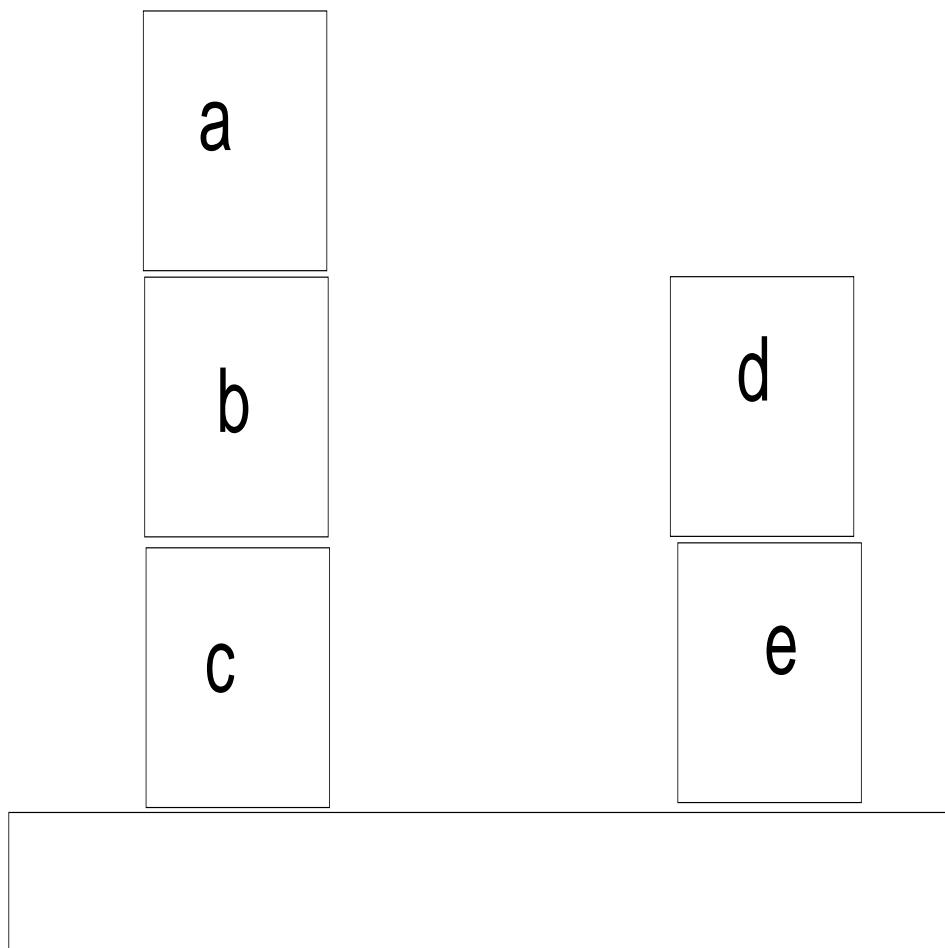


Blocks World Example



Domain and Relations

domain

$$\mathcal{D} = \{a, b, c, d, e\}$$

relations

$$\{\langle a, b \rangle, \langle b, c \rangle, \langle d, e \rangle\}$$

$$\{\langle a, b \rangle, \langle b, c \rangle, \langle a, c \rangle, \langle d, e \rangle\}$$

$$\{c, e\}$$

$$\{a, d\}$$

Functions

functions

$$\{\langle b \rangle \rightarrow a, \langle c \rangle \rightarrow b, \langle e \rangle \rightarrow d\}$$

Non-Logical Symbols

Predicate Symbols

ABOVE

TABLE

CLEAR

ON

Function Symbols

BLOCKA

BLOCKB

BLOCKC

BLOCKD

BLOCKE

Interpretation Mapping

$$\mathcal{I}[\text{ON}] = \{\langle a, b \rangle, \langle b, c \rangle, \langle d, e \rangle\}$$

$$\mathcal{I}[\text{ABOVE}] = \{\langle a, b \rangle, \langle b, c \rangle, \langle a, c \rangle, \langle d, e \rangle\}$$

$$\mathcal{I}[\text{BLOCKA}] = a$$

$$\mathcal{I}[\text{BLOCKB}] = b$$

$$\mathcal{I}[\text{BLOCKC}] = c$$

$$\mathcal{I}[\text{BLOCKD}] = d$$

$$\mathcal{I}[\text{BLOCKE}] = e$$

$$\mathcal{I}[\text{HAT}] = \{\langle b \rangle \rightarrow a, \langle c \rangle \rightarrow b, \langle e \rangle \rightarrow d\}$$

$$\mathcal{I}[\text{TABLE}] = \{c, e\}$$

$$\mathcal{I}[\text{CLEAR}] = \{a, d\}$$

Another Interpretation (Non-Intended)

$$\mathcal{I}[\text{ABOVE}] = \{\langle b, a \rangle, \langle c, b \rangle, \langle e, d \rangle\}$$

$$\mathcal{I}[\text{ON}] = \{\langle b, a \rangle, \langle c, b \rangle, \langle c, a \rangle, \langle e, d \rangle\}$$

$$\mathcal{I}[\text{BLOCKA}] = a$$

$$\mathcal{I}[\text{BLOCKB}] = c$$

$$\mathcal{I}[\text{BLOCKC}] = b$$

$$\mathcal{I}[\text{BLOCKD}] = d$$

$$\mathcal{I}[\text{BLOCKE}] = e$$

$$\mathcal{I}[\text{HAT}] = \{\langle b \rangle \rightarrow a, \langle c \rangle \rightarrow b, \langle e \rangle \rightarrow d\}$$

$$\mathcal{I}[\text{CLEAR}] = \{c, e\}$$

$$\mathcal{I}[\text{TABLE}] = \{a, d\}$$

Example

Does the following formula hold in the intended interpretation?

$$\mathfrak{I} \models \text{ABOVE}(\text{BLOCKA}, \text{BLOCKB})$$

Example

Does the following formula hold in the intended interpretation?

$$\mathfrak{S} \models \text{ABOVE}(\text{HAT(BLOCKB)}, \text{BLOCKB})$$

Example

Does the following formula hold in the intended interpretation?

$$\mathfrak{S} \models \text{ABOVE}(\text{BLOCKA}, \text{BLOCKB}) \wedge \text{ABOVE}(\text{BLOCKA}, \text{BLOCKC})$$

Variables

$\text{ON}(x, y) \rightarrow \text{ABOVE}(x, y)$

Variable Assignment μ

$$\mu(x) = \text{a}$$

$$\mu(y) = \text{a}$$

$$\mu(z) = \text{b}$$

Example

$\mathfrak{I} \mu' \models \forall x.\text{ON}(x, \text{BLOCKA})$

Example

$$\forall x, y \text{ON}(x, y) \rightarrow \text{ABOVE}(x, y)$$

Example

$\exists y \text{ ABOVE}(y, \text{BLOCK})$

Denotation

variable assignment μ over \mathcal{D}

1. If x is a variable then $\llbracket x \rrbracket_{\mathfrak{S}, \mu} = \mu[x]$
2. If t_1, \dots, t_n are terms, and F is a function symbol of arity n then

$$\llbracket F(t_1, \dots, t_n) \rrbracket_{\mathfrak{S}, \mu} = \mathcal{I}(F)(\llbracket t_1 \rrbracket_{\mathfrak{S}, \mu}, \dots, \llbracket t_n \rrbracket_{\mathfrak{S}, \mu})$$

Satisfaction

1. $\mathfrak{S}, \mu \models P(t_1 \dots t_n)$ iff $\langle d_1, \dots, d_n \rangle \in \mathcal{I}(P)$,
where $d_i = \llbracket t_i \rrbracket_{\mathfrak{S}, \mu}$
2. $\mathfrak{S}, \mu \models t_1 = t_2$ iff $\llbracket t_1 \rrbracket_{\mathfrak{S}, \mu}$ and $\llbracket t_2 \rrbracket_{\mathfrak{S}, \mu}$ and are
the same element of D.
3. $\mathfrak{S}, \mu \models \neg \alpha$ iff it is not the case that $\mathfrak{S}, \mu \models \alpha$
4. $\mathfrak{S}, \mu \models (\alpha \wedge \beta)$ iff $\mathfrak{S}, \mu \models \alpha$ and
 $\mathfrak{S}, \mu \models \beta$.
5. $\mathfrak{S}, \mu \models (\alpha \vee \beta)$ iff $\mathfrak{S}, \mu \models \alpha$ or
 $\mathfrak{S}, \mu \models \beta$.
6. $\mathfrak{S}, \mu \models \exists x. \alpha$ iff $\mathfrak{S}, \mu' \models \alpha$ for some variable
assignment μ' that differs from μ on at most
 x .
7. $\mathfrak{S}, \mu \models \forall x. \alpha$ iff $\mathfrak{S}, \mu' \models \alpha$ for every variable
assignment μ' that differrs from μ on at most
 x .

Another Example

$\mathcal{D} =$

{GeorgeBush, ArnoldSchwartzenegger,
BillClinton, MargaretThatcher,
Aristotle, Mozart}

Non-Logical Symbols

Predicate Symbols

H

M

LOVES

Function Symbols

ARNIE

BILL

GEORGE

MAGGIE

ARI

WOLFGANG

Interpretation Mapping

$\mathcal{I}[\text{LOVES}] =$

$\{\langle \text{ArnoldSchwartzenegger}, \text{ArnoldSchwartzenegger} \rangle$
 $\langle \text{GeorgeBush}, \text{MargaretThatcher} \rangle,$
 $\langle \text{MargaretThatcher}, \text{GeorgeBush} \rangle,$
 $\langle \text{Mozart}, \text{BillClinton} \rangle,$
 $\langle \text{BillClinton}, \text{Aristotle} \rangle,$
 $\langle \text{Aristotle}, \text{Mozart} \rangle\}$

$\mathcal{I}[\text{ARNIE}] = \text{ArnoldSchwartzenegger}$

$\mathcal{I}[\text{BILL}] = \text{BillClinton}$

$\mathcal{I}[\text{GEORGE}] = \text{GeorgeBush}$

$\mathcal{I}[\text{MAGGIE}] = \text{MargaretThatcher}$

$\mathcal{I}[\text{WOLFGANG}] = \text{Mozart}$

$\mathcal{I}[\text{ARI}] = \text{Aristotle}$

Interpretation Mapping(cont)

$$\mathcal{I}(H) = \{GeorgeBush, ArnoldSchwartzenegger, BillClinton, MargaretThatcher, Aristotle, Mozart\}$$
$$\mathcal{I}(M) = \{GeorgeBush, MargaretThatcher\}$$

Statements to Check

$H(MAGGIE)$

$M(WOLFGANG)$

$L(WOLFGANG, ARI)$

$L(ARNIE, ARNIE)$

$\forall x H(x)$

$\exists x M(x)$

Statements to Check

$M(ARNIE) \rightarrow \exists y L(GEORGE, y)$

$M(MAGGIE) \rightarrow \exists y L(GEORGE, y)$

Statements to Check

$$\exists x H(x) \wedge M(x)$$